

Recap

Midterm next week!

$$\text{SVD: } A = \sum_{i=1}^n \sigma_i w_i v_i^*$$

$$\text{Rank-}k \text{ approximation: } A_k = \sum_{i=1}^k \sigma_i w_i v_i^*$$

$$\forall B, \text{rank}(B) \leq k, \|A - B\|_2 \geq \sigma_{k+1} = \|A - A_k\|_2$$

Least squares approximation: $a_1, \dots, a_n \in \mathbb{R}^d$, $\dim(S) \leq k$
(PCA)

$$\sum_{i=1}^n \text{dist}(a_i, S) \geq \sigma_1^2 + \dots + \sigma_k^2$$

$$A = \begin{bmatrix} \text{---} & a_1^T & \text{---} \\ \vdots & \vdots & \vdots \\ \text{---} & a_n^T & \text{---} \end{bmatrix} \quad \text{Opt} = \text{Span of top } k \text{ right singular vectors}$$

Gershgorin Disc Theorem

Given $M \in \mathbb{C}^{n \times n}$

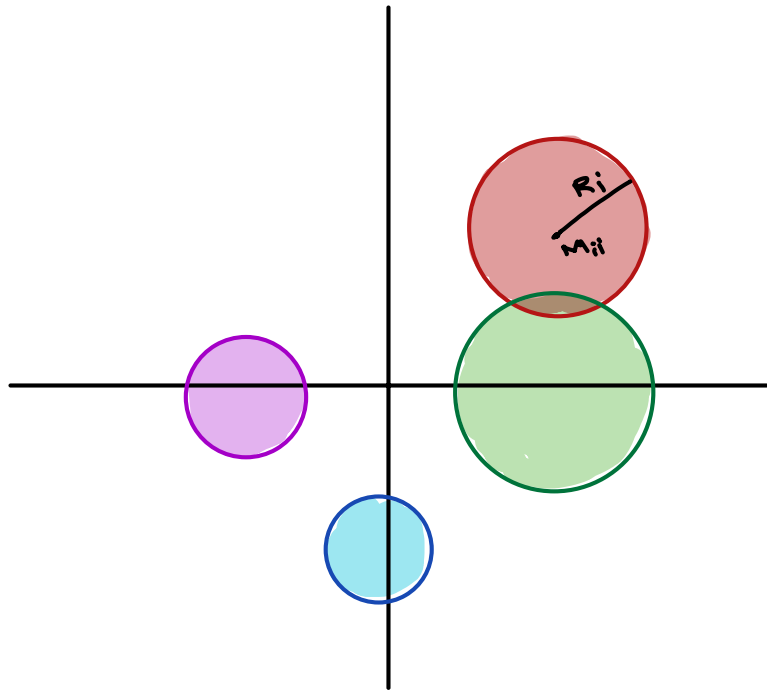
$$M = \begin{pmatrix} a_1 & & & \\ & a_2 & & \\ & & \ddots & \\ & & & a_n \end{pmatrix}$$

eigenvalues of M are $a_1, \dots, a_n$

What about general $M \in \mathbb{C}^{n \times n}$

► Let $M \in \mathbb{C}^{n \times n}$. Let $R_i = \sum_{j \neq i} |M_{ij}|$ and $\text{Disc}(M_{ii}, R_i) = \{z \mid |z - M_{ii}| < R_i\}$

Then, for any eigenvalue λ of M , $\lambda \in \bigcup_i \text{Disc}(M_{ii}, R_i)$



Proof: Let x be an eigenvector with eigenvalue $\lambda \in \mathbb{C}$.

Let $i_0 = \arg \max_{i \in [n]} |x_i|$

$$Mx = \lambda x \Rightarrow (Mx)_{i_0} = \lambda x_{i_0}$$

$$\Rightarrow \sum M_{i_0 j} x_j = \lambda x_{i_0}$$

$$\Rightarrow M_{i_0 i_0} + \sum_{j \neq i_0} M_{i_0 j} \frac{x_j}{x_{i_0}} = \lambda$$

$$\begin{aligned} \therefore |\lambda - M_{i_0 i_0}| &\leq \sum_{j \neq i_0} |M_{i_0 j}| \cdot \underbrace{\frac{|x_j|}{|x_{i_0}|}}_{\leq 1} \\ &\leq R_{i_0} \end{aligned}$$

An application to Compressed Sensing

- ▶ $A \in \mathbb{R}^{k \times n}$ is said to have **restricted isometry property** with parameters (S, δ) if

$$(1-\delta) \|x\|^2 \leq \|Ax\|^2 \leq (1+\delta) \|x\|^2$$

for all S -sparse vectors x s.t. $|\{i \mid x_i \neq 0\}| \leq S$

↪ useful condition, NP-hard to check.

- ▶ $A \in \mathbb{R}^{k \times n}$

$$A = \begin{bmatrix} | & & & | \\ a_1 & \dots & & a_n \\ | & & & | \end{bmatrix} \quad \text{with } \|a_i\| = 1 \quad \forall i$$

Coherence $\mu(A) = \max_{i \neq j} |\langle a_i, a_j \rangle|$

Incoherence implies restricted isometry

For $A \in \mathbb{R}^{k \times n}$ as before with coherence $\mu(A)$,

A has restricted isometry property with parameters $(S, (S-1) \cdot \mu(A))$ i.e. $\forall x$ s.t. $|\{i \mid x_i \neq 0\}| \leq S$

$$(1 - (S-1) \cdot \mu(A)) \cdot \|x\|^2 \leq \|Ax\|^2 \leq (1 + \overbrace{(S-1) \cdot \mu(A)}^{\delta}) \cdot \|x\|^2$$

Proof: Say $S=1$ $\|Ax\|^2 = x_i^2 = \|x\|^2$

$\text{If } S = \{i \mid x_i \neq 0\}$ $\|Ax\|^2 = \|A_S x_S\|^2$ only keep columns for $i \in S$

$$= \langle x_S, A_S^* A_S x_S \rangle$$

$$A_S^* A_S = \begin{bmatrix} 1 & & \\ & \ddots & \\ & & 1 \end{bmatrix}$$

$|\langle a_i, a_j \rangle| \leq \mu(A)$

$\therefore$ by Gershgorin, $|\lambda - 1| \leq (S-1)\mu(A) = \delta$
 $\forall$ eigenvalues λ

$$\frac{\langle x_S, A_S^* A_S x_S \rangle}{\|x_S\|^2} \in [1 - \delta, 1 + \delta]$$

Solving linear equations

$$Ax = b$$

$$\begin{cases} x + y + z = 10 \\ x + 2y + 3z = 15 \\ 2x + 3y + z = 19 \end{cases}$$

$$\begin{aligned} x + y + z &= 10 \\ y + 2z &= 5 \\ y - z &= -1 \end{aligned}$$

$$\begin{aligned} x &= 7 \\ y &= 1 \\ z &= 2 \end{aligned}$$
$$\begin{aligned} x + y + z &= 10 \\ y + 2z &= 5 \\ -3z &= -6 \end{aligned}$$

$$[A \mid b]$$

$$\begin{bmatrix} 1 & 1 & 1 & 10 \\ 1 & 2 & 3 & 15 \\ 2 & 3 & 1 & 19 \end{bmatrix}$$

$$\begin{aligned} &\downarrow \\ &\left(\begin{array}{ccc|c} 1 & 1 & 1 & 10 \\ & 1 & 2 & 5 \\ & 1 & -1 & -1 \end{array} \right) \\ &\downarrow \\ &\vdots \end{aligned}$$

Row-reduced form

$$M = \begin{bmatrix} 1 & 1 & 1 & 10 \\ 0 & 1 & 2 & 15 \\ 0 & 0 & 1 & 2 \end{bmatrix}$$

- First non-zero entry of each row is 1. (leading entry)
- If M_{ij} is leading entry then M is 0 "below" & "left"
($M_{i'j'} = 0 \forall i' \geq i, j' \geq j$)
- Rows containing all 0s at bottom.

$$\begin{bmatrix} 1 & 1 & 1 & 10 \\ 1 & 2 & 3 & 15 \\ 2 & 3 & 1 & 19 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 1 & 1 & 10 \\ 0 & 1 & 2 & 5 \\ 0 & 1 & -1 & -1 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 1 & 1 & 10 \\ 0 & 1 & 2 & 5 \\ 0 & 0 & -3 & -6 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 1 & 1 & 10 \\ 0 & 1 & 2 & 5 \\ 0 & 0 & 1 & 2 \end{bmatrix}$$

Elementary row-operations

- $M_i \leftarrow M_i + c \cdot M_j$
- $M_i \leftarrow c \cdot M_i \quad (c \neq 0)$
- Swap M_i, M_j

Properties

Ex: Matrix in row-reduced form "obviously" has
row-rank = column-rank

e.g.
$$\begin{bmatrix} 1 & 1 & 1 & 10 \\ 0 & 1 & 2 & 15 \\ 0 & 0 & 1 & 2 \end{bmatrix}$$

Ex: Elementary row-operations do not change
row-rank or column-rank

Ex: Above is true over any field $\mathbb{F}$.

► For any matrix $M \in \mathbb{F}^{m \times n}$, row-rank = column-rank

Solving (sparse) linear systems fast

Given $A_0 x = b_0$ $\longrightarrow$ $\underbrace{A_0^T A_0}_A x = \underbrace{A_0^T b_0}_b$

(say s non-zero entries in A_0)

$$A x = b$$

$$\ker(A_0) = \{0\}$$

$$\iff$$

A is positive definite

$$\iff$$

$$\exists x.$$

$\exists$ unique solution, say u

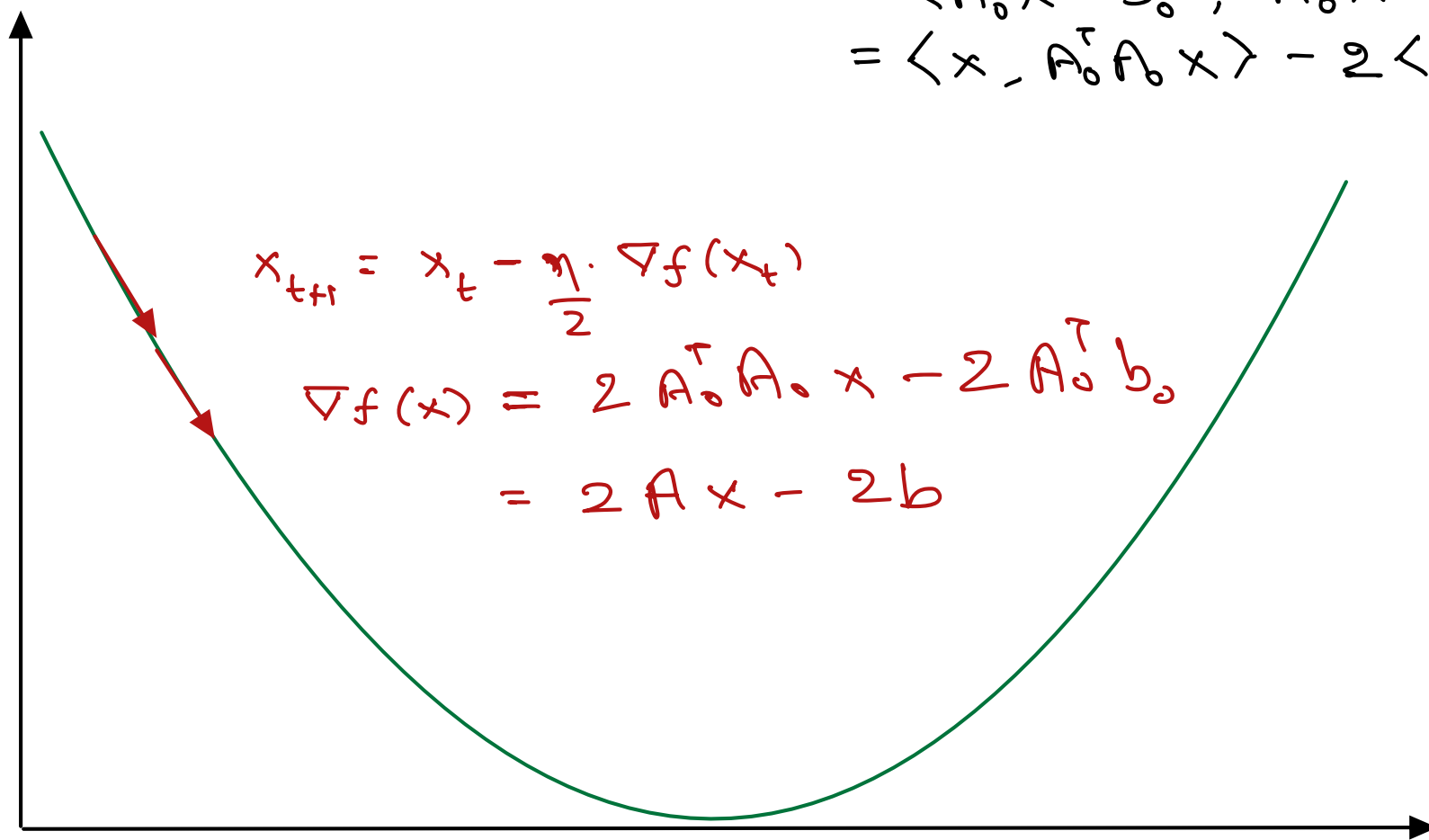
Goal: Compute approximate solution fast.

Gold standard: Time to check solution
= matrix-vector product
= ?

Solves from optimization

Consider: minimize $f(x) = \|A_0 x - b_0\|^2$

$$\begin{aligned} &= \langle A_0 x - b_0, A_0 x - b_0 \rangle \\ &= \langle x, A_0^T A_0 x \rangle - 2 \langle x, A_0^T b_0 \rangle + \|b_0\|^2 \end{aligned}$$



$$x_{t+1} = x_t - \frac{\eta}{2} \nabla f(x_t)$$

$$\begin{aligned} \nabla f(x) &= 2 A_0^T A_0 x - 2 A_0^T b_0 \\ &= 2 A x - 2 b \end{aligned}$$

$$x_{t+1} = x_t - \eta (A x_t - b)$$

Steepest descent iteration

► Let u be the solution to $Ax=b$. Then

$$(x_t - u) = (I - \eta A)^t (x_0 - u)$$

Proof:

$$x_{t+1} = x_t - \eta (Ax_t - Au)$$

$$x_{t+1} - u = x_t - u - \eta A (x_t - u)$$

$$x_{t+1} - u = (I - \eta A) (x_t - u)$$

Convergence analysis : choosing η

$$(x_t - u) = (I - \eta A)^t (x_0 - u) \Rightarrow \|x_t - u\| \leq \|(I - \eta A)^t\|_2 \cdot \|x_0 - u\|$$

Eigenvalues

$$A : \lambda_1 \geq \dots \geq \lambda_n > 0 \text{ (say)}$$

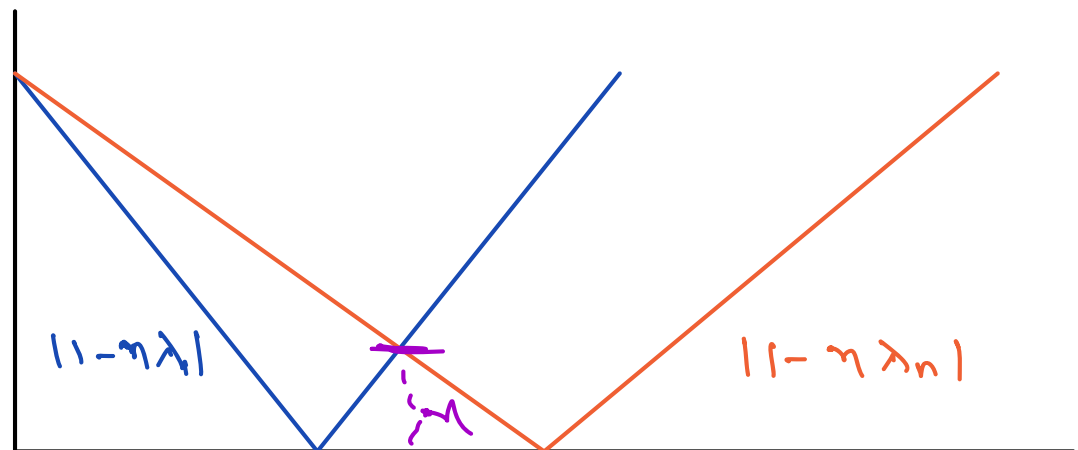
$$I - \eta A : 1 - \eta \lambda_1 \leq \dots \leq 1 - \eta \lambda_n < 1$$

$$(I - \eta A)^t : (1 - \eta \lambda_1)^t, \dots, (1 - \eta \lambda_n)^t$$

$$\|(I - \eta A)^t\|_2 = \max \{ |1 - \eta \lambda_1|^t, |1 - \eta \lambda_n|^t \}$$

$$\eta \lambda_1 - 1 = 1 - \eta \lambda_n$$

$$\eta = \frac{2}{\lambda_1 + \lambda_n}$$



Convergence analysis

$$\|(\mathbf{I} - \eta \mathbf{A})^t\|_2 = (1 - \eta \lambda_n)^t = \left(1 - \frac{2 \lambda_n}{\lambda_1 + \lambda_n}\right)^t = \left(1 - \frac{2}{\underbrace{(\lambda_1/\lambda_n) + 1}_{\kappa}}\right)^t$$

Condition number

$$\|(\mathbf{I} - \eta \mathbf{A})^t\|_2 \leq \epsilon \quad \text{after} \quad t = O\left(\kappa \log \frac{1}{\epsilon}\right)$$

Another distance: $\|x_t - u\|_A = \sqrt{\underbrace{(x_t - u)^T A (x_t - u)}_{\text{inner product?}}}$

Ex: Show that same convergence analysis also works for $\|x - u\|_A$